Optimal decision trees using optimization techniques MASTER THESIS Master in Data Science

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July 8, 2019



• Optimal decision trees

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• Optimal decision trees



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• Optimal decision trees

2 Our solution

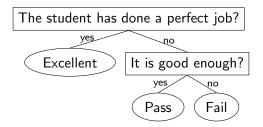


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Project's introduction

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Optimal decision tree iif:

- 100% accuracy on train
- 2 With a smaller size \implies accuracy \neq 100%

Proof

We need proof that this specific tree is optimal

$(\neg v_1)$

$$\begin{pmatrix} \neg v_1 \end{pmatrix} \\ v_i \to \neg l_{ij} \qquad j \in \mathrm{LR}(i)$$

$$\begin{pmatrix} \neg v_1 \end{pmatrix} \\ v_i \to \neg l_{ii} & j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} & j \in \mathrm{LR}(i) \end{cases}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i,i} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg v_i \rightarrow \left(\sum_{j \in \mathrm{LR}(i)} l_{ij} = 1\right) \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{ii} & j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} & \neg v_i \rightarrow \left(\sum_{i \in \mathrm{LR}(i)} l_{ij} = 1 \right) \\ p_{ji} \leftrightarrow l_{ij}, & j \in \mathrm{LR}(i) \\ p_{ji} \leftrightarrow r_{ij}, & j \in \mathrm{RR}(i) \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{ii} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg v_i \rightarrow \left(\sum_{\substack{i \in \mathrm{LR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \quad i \in \mathrm{RR}(i) \\ \begin{pmatrix} \min(-1,N) \\ i \in j \\ i \in j \end{pmatrix}} v_{ij} \right) \\ \text{with } j = 2, \dots, N \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \to \neg l_{i:i} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg v_i \to \left(\sum_{\substack{i \in \mathrm{RR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \qquad i \in \mathrm{RR}(i) \\ \vdots & \vdots & (\sum_{\substack{i \in \mathrm{RR}(i) \\ \sum p_{ji} \to (r_{ij}, d_{r_i}^{j_i}) \vee (a_{r_i} \wedge r_{ij})) \\ \vdots & \vdots & (\sum_{\substack{i \in \mathrm{IRR}(i) \\ \sum p_{i=1} \\ i \neq i} ((p_{ji} \wedge d_{r_i}^{j_i}) \vee (a_{r_i} \wedge r_{ij}))); d_{r_i}^n = 0. \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i:} \quad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \quad \neg v_i \rightarrow \left(\sum_{\substack{i \in \mathrm{LR}(i) \\ j \in \mathrm{LR}(i)}} l_{ij} = 1 \right) \\ p_{ji} \leftrightarrow l_{ij}, \quad j \in \mathrm{LR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \quad \substack{i \in \mathrm{RR}(i) \\ \sum_{\substack{(\min(j-1,N) \\ \sum} p_{ji} = 1)}} \\ d_{r_j}^{\flat} \leftrightarrow \left(\sum_{\substack{i=1 \atop i \in \mathrm{I}}}^{j-1} ((p_i \wedge d_{r_i}^i) \vee (a_{r_i} \wedge r_{ij})) \right); d_{r_i}^{\flat} = 0. \\ d_{r_j}^{\flat} \leftrightarrow \left(\sum_{\substack{i=1 \atop i \in \mathrm{I}}}^{j-1} ((p_i \wedge d_{r_i}^i) \vee (a_{r_i} \wedge l_{ij})) \right); d_{r_i}^{\flat} = 0. \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i:i} \quad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \quad \neg v_i \rightarrow \left(\sum_{\substack{i \in \mathrm{LR}(i) \\ j \in \mathrm{V}_i(j)}} l_{ij} = 1 \right) \\ p_{ji} \leftrightarrow r_{ij}, \quad j \in \mathrm{LR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \quad i \in \mathrm{RR}(i) \\ (\sum_{\substack{i \in \lfloor j \\ i = \lfloor j \end{bmatrix}} p_{ji} = 1) \quad \text{with } j = 2, \dots, N \\ d_{r_j}^{i} \leftrightarrow \left(\sum_{\substack{i \in \lfloor j \\ i = \lfloor j \end{bmatrix}} (p_{ji} \rightarrow d_{r_i}) \vee (a_{r_i} \wedge r_{ij}) \right); d_{r_i}^{i} = 0. \\ \sum_{\substack{i = \lfloor j \\ i = \lfloor j \end{bmatrix}} (u_{r_i} \wedge p_{ji} \rightarrow \neg a_{r_j}) \quad \stackrel{l_i) \vee (a_{r_i} \wedge l_{ij})}{(u_r \rightarrow p_{ji})}); d_{r_i}^{i} = 0.$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i:} \quad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \quad \neg v_i \rightarrow \left(\sum_{\substack{i \in \mathrm{LR}(i) \\ j \in \mathrm{LR}(i)}} l_{ij} = 1 \right) \\ p_{ji} \leftrightarrow l_{ij}, \quad j \in \mathrm{LR}(i) \\ (\sum_{\substack{min(j-1,N) \\ (min(j-1,N) \\ j \in \mathrm{LR}(i)}} l_{ij} = 1) \\ d_{r_j}^j \leftrightarrow \left(\sum_{\substack{i \in \mathrm{II} \\ i \in \mathrm{II}}} l_{ij} (p_{ji} \wedge d_{r_i}^0) \vee (a_{r_i} \wedge r_{ij}) \right); d_{r_i}^1 = 0. \\ (\sum_{\substack{i \in \mathrm{II} \\ i \in \mathrm{II}}} l_{ij} \vee (a_{r_i} \wedge p_{ji} \rightarrow \neg a_{r_j}) \\ u_{r_j} \leftrightarrow \left(a_{r_j} \vee \sum_{\substack{i=1 \neq \mathrm{II} \\ i=1 \neq \mathrm{II}}} l_{ij} \vee (u_{r_i} \wedge p_{ji}) \right) \\ \neg v_j \rightarrow \left(\sum_{r=1}^{K} a_{r_j} = 1 \right) \\ \end{array} \right) \text{ with } j = 1, \dots, N$$

$$\begin{array}{c} (\neg \upsilon_1) \\ \upsilon_i \to \neg l_{i:i} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg \upsilon_i \to \left(\sum_{\substack{i \in \mathrm{RR}(i) \\ j \in \mathrm{RR}(i)}} l_{ij} = 1 \right) \\ p_{ji} \leftrightarrow l_{ij}, \qquad j \in \mathrm{LR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \qquad i \in \mathrm{RR}(i) \\ \begin{pmatrix} \min(j^{-1,N)} \\ p_{ji} = 1 \end{pmatrix} \qquad \text{with } j = 2, \dots, N \\ \begin{pmatrix} d_{r_j} \leftrightarrow \left(\sum_{i=1}^{j-1} (u_{r_i} \wedge r_{ij}) \right) \\ i \in \lfloor \frac{j}{2} \end{bmatrix}} \\ \ell^{j-1} \qquad (u_{r_i} \wedge p_{ji} \to a_{r_j}) \\ u_{r_j} \leftrightarrow \left(a_{r_j} \vee \sum_{i=1}^{j-1} (u_{r_i} \wedge p_{ji}) \\ \neg \upsilon_j \to \left(\sum_{i=1}^{K} a_{r_j} = 1 \right) \\ \neg \upsilon_j \to \left(\sum_{r=1}^{K} a_{r_j} = 0 \right) \qquad \text{with } j = 1, \dots, N \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i:i} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg v_i \rightarrow \left(\sum_{\substack{i \in \mathrm{LR}(i) \\ j \in \mathrm{CR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \qquad i \in \mathrm{RR}(i) \\ \dots \qquad (\sum_{\substack{(j) \in \mathbb{N}^{(j)} \\ \sum_{i \in [j]} (v_{ii} \wedge d_{ii}^{n}) \vee (a_{ii} \wedge r_{ij})) \\ \sum_{i \in [j]} (u_{ii} \wedge p_{ji} \rightarrow \neg a_{rj}) \qquad \text{with } j = 2, \dots, N \\ \frac{d_{ij} \leftrightarrow \left(\sum_{i=1}^{j-1} ((v_{ii} \wedge p_{ji}) \vee (a_{ii} \wedge r_{ij})) \right); d_{i,1}^{i} = 0. \\ \dots \qquad (\sum_{\substack{(j) \in \mathbb{N}^{(j)} \\ i = \lfloor j \rfloor \\ u_{rj} \leftrightarrow \left(a_{rj} \vee \sum_{i=1}^{j-1} (u_{ri} \wedge p_{ji}) \right) \\ \neg v_{j} \rightarrow \left(\sum_{i=1}^{K} a_{rj} = 1 \right) \qquad \text{with } j = 1, \dots, N \\ v_{j} \rightarrow \left(\sum_{r=1}^{K} a_{rj} = 0 \right) \qquad v_{j} \wedge \neg c_{j} \rightarrow \bigvee_{r=1}^{K} d_{r,j}^{\sigma}(r,q) \\ \end{array}$$

$$\begin{array}{c} (\neg v_1) \\ v_i \rightarrow \neg l_{i:} \qquad j \in \mathrm{LR}(i) \\ l_{ij} \leftrightarrow r_{ij+1} \qquad \neg v_i \rightarrow \left(\sum_{\substack{\mathrm{LR}(i) \\ j \in \mathrm{RR}(i) \\ p_{ji} \leftrightarrow r_{ij}, \qquad i \in \mathrm{RR}(i) \\ (\sum_{\substack{(i) \in (j_i) \land d_i > (\sum_{i \in \mathrm{RR}(i) \\ j \in \mathrm{RR}(i) \\ (\sum_{i \in [j]} (u_{i} \land p_{ji}) \rightarrow (a_{i} \land r_{ij})) \\ i \in [j] \qquad (i_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (j \neq (\sum_{i \in [j]} (u_{i} \land p_{ji})); d_{i,1}^{\theta} = 0. \\ (\sum_{\substack{(i) \in [j] \\ i \in [j] \\ i \in [j] \\ (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (j \neq (a_{i} \land l_{ij})) \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ji} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{i} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{ij} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (a_{i} \land l_{ij})); d_{i,1}^{\theta} = 0. \\ (\sum_{i \in [j]} (u_{i} \land p_{i} \rightarrow \neg a_{rj}) \qquad (i_{i}) \lor (i_{i} \land i_{i} \land i_{i} \rightarrow \neg a_{rj}) \land (i_{i}) \lor (i_{i} \land i_{i} \land i_{i} \rightarrow \neg a_{rj}) \land (i_{i}) \lor (i_{i} \land i_{i} \rightarrow \neg a_{rj} \land (i_{i} \land i_{rj} \land (i_{i} \land i_{rj} \land (i_{i} \land i_{rj} \land i_{rj} \land (i_{i} \land i_{rj} \land (i_{i} \land i_{rj} \land i_{rj} \land (i_{i} \land i_{rj} \land (i_{i} \land i_{rj} \land i_{rj} \land (i_{i} \land i_{rj} \land (i_{rj} \land i_{rj} \land (i_{i}$$

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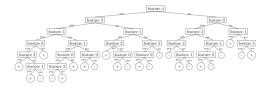
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Problem

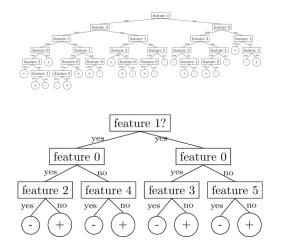
Obtaining optimal decision trees is hard

Is it really a problem, though? Why would an optimal decision tree interest me? Sklearn is good enough, right?



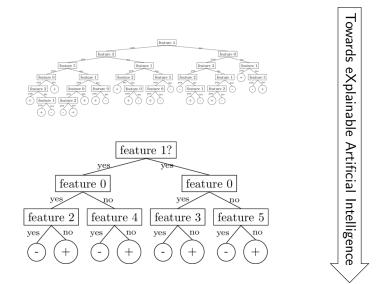
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We have tools for greedy algorithms

We have tools for greedy algorithms

We don't have tools for optimal trees

Software package that makes optimal algorithms user friendly Prepare the data Software package that makes optimal algorithms user friendly

- Prepare the data
- Pind optimal trees

Software package that makes optimal algorithms user friendly

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Software package that makes optimal algorithms user friendly

- Prepare the data
- Pind optimal trees
- Solution API Abstract different CP solvers and algorithms into a common API
- Return the solutions as a ready-to-use class

Convert this...

Algorithm 1: Searching for the optimal decision tree

input: A dataset *data* as a matrix of $n_{feats} + 1 \times n_{rows}$

 $search_space \leftarrow SOMETHING;$ //specific of each encoding

while more_params_available(search_space) do

```
params ← next_params(search_space);
model ← encode(data, params);
solution ← solve(model);
if solution is optimal then
  | return solution
end
```

end

return "No solution"

```
...into this:
    model = SomeDTModel(...)
    data = np.array(...)
    tree = model.find_tree(data)
```

tree is proven to be optimal!

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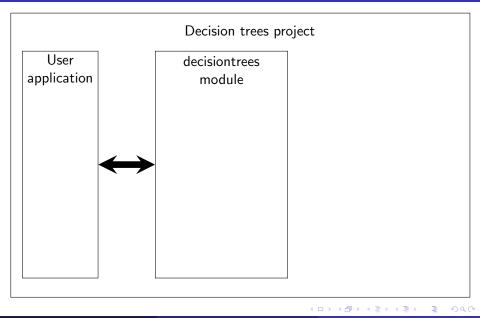
Implemented models:

- Optimize depth: Avellaneda, 2020
- Optimize size (using all observations)
 - 2x Narodytska et al, 2018
 - 1x Avellaneda, 2020
- Optimize size (using a subset of observations): Avellaneda, 2020
- Bonus, non-optimal Scikit-Learn DecisionTreeClassifier

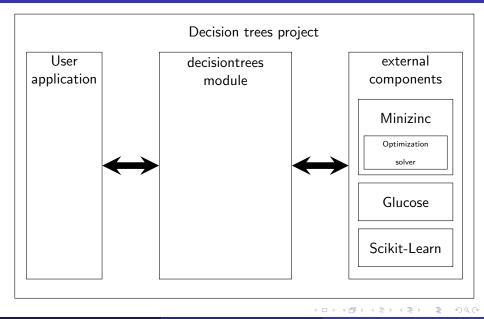
Our solution - Overview



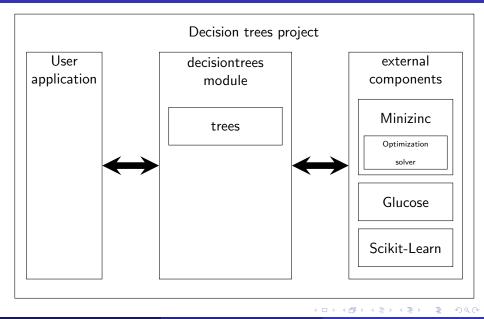
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Our solution - Overview

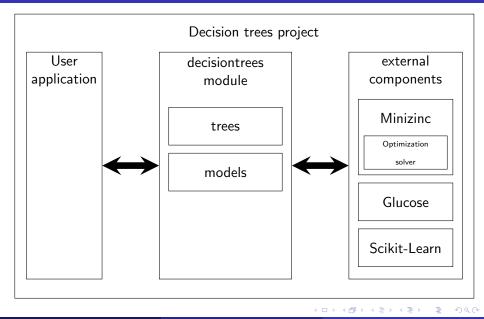


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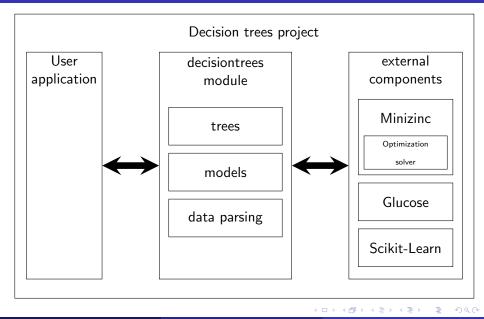
Optimal decision trees

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Our solution - Overview



Our solution - Overview



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Optimal decision trees

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• Mathematical optimization approaches: \uparrow smaller trees; \downarrow slow

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- Greedy approaches: \uparrow fast; \downarrow don't find (nor proof) optimal trees

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- Slowness \leftarrow lack of interfaces between Minizinc and SAT solvers

- Mathematical optimization approaches: \uparrow smaller trees; \downarrow slow
- Greedy approaches: \uparrow fast; \downarrow don't find (nor proof) optimal trees
- Slowness \leftarrow lack of interfaces between **Minizinc** and SAT solvers **POTENTIAL FUTURE WORK!**

• Interface Minizinc with SAT solvers

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- Interface Minizinc with SAT solvers
- Add more approaches!!!

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- Interface Minizinc with SAT solvers
- Add more approaches!!!
- (Neverending future work) Improve source code!

- Towards XAI: Optimal decision trees are smaller
- Created a successful tool to find optimal decision trees
- Breaking entry barriers for non-technical users

Thanks for watching